
cREXX Progress Update

The 33rd Annual Rexx Symposium

Adrian Sutherland • 12.09.2022 (Final)

cREXX Progress Update

cREXX Architecture

cREXX Level B MVP

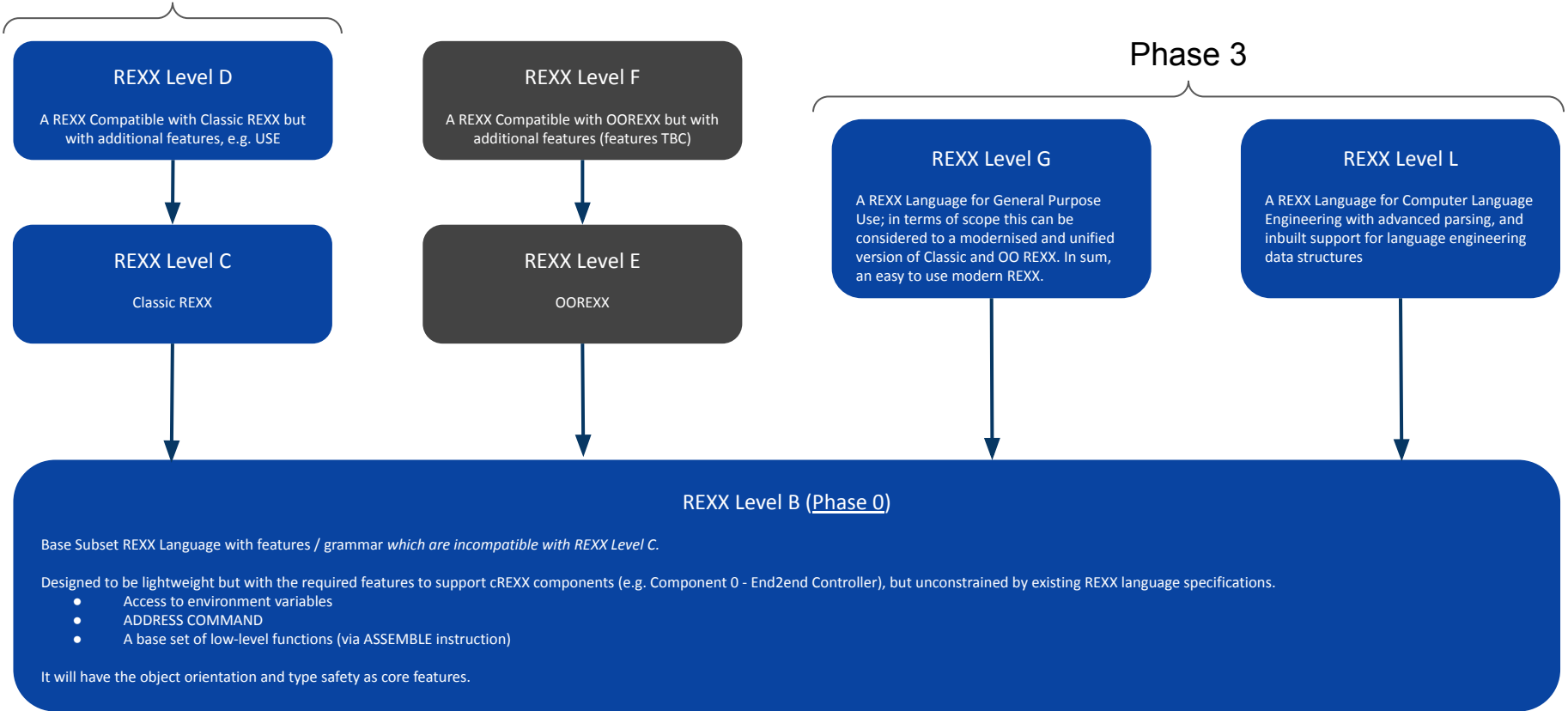
How to Help?

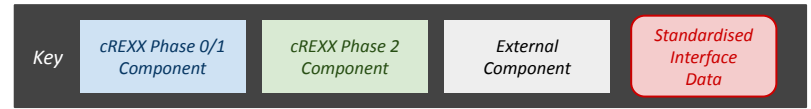
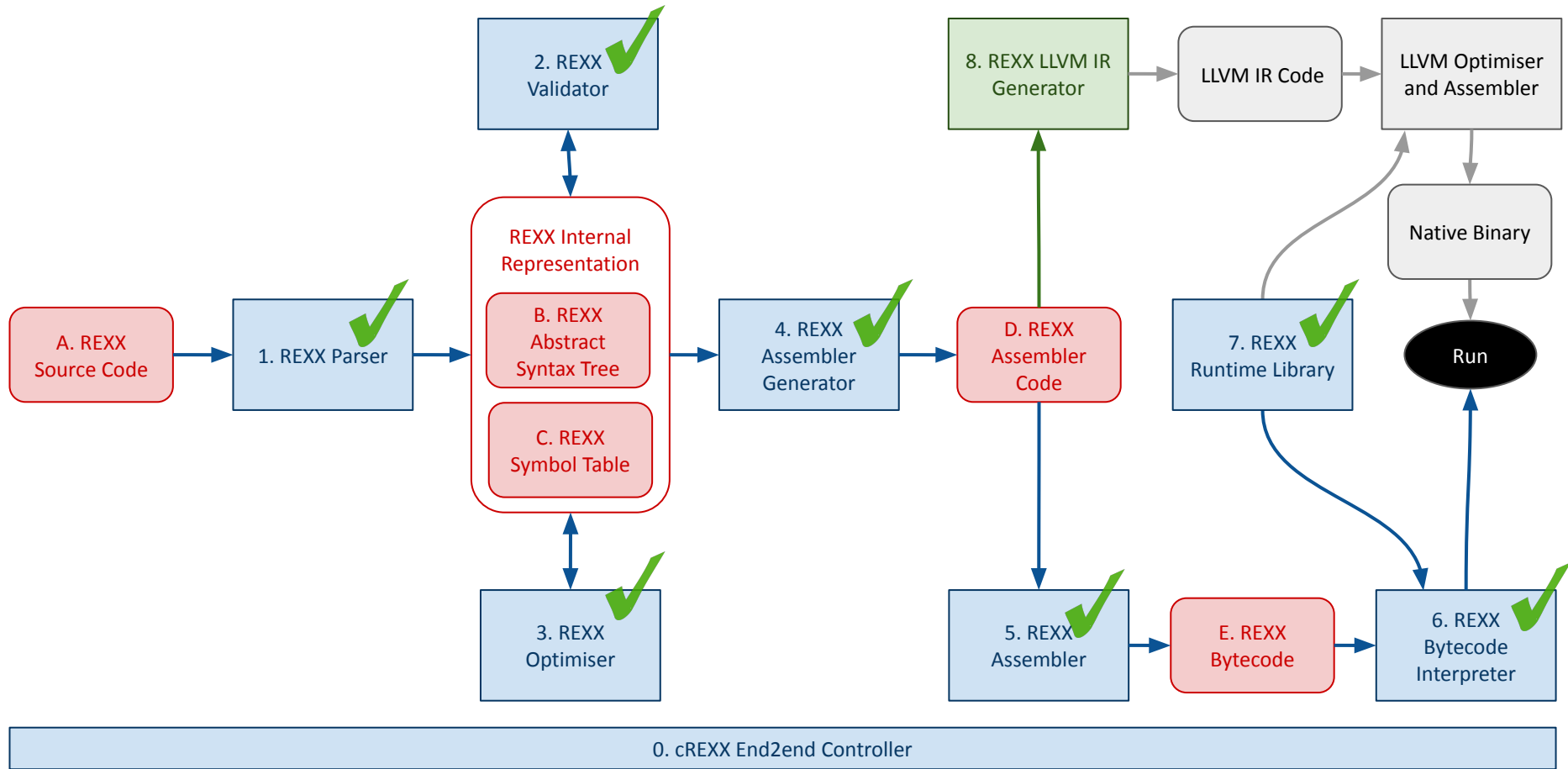
Thanks!

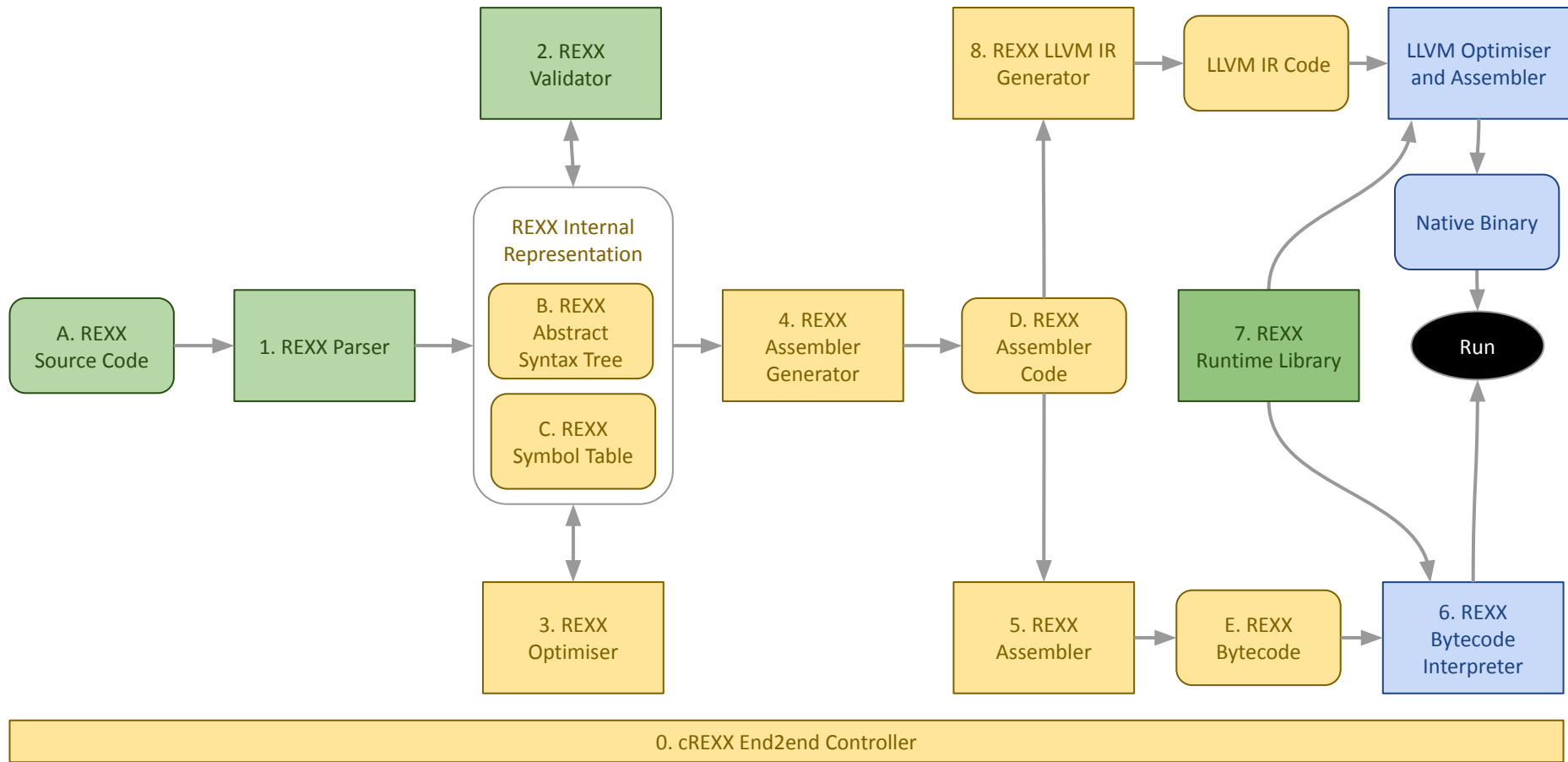


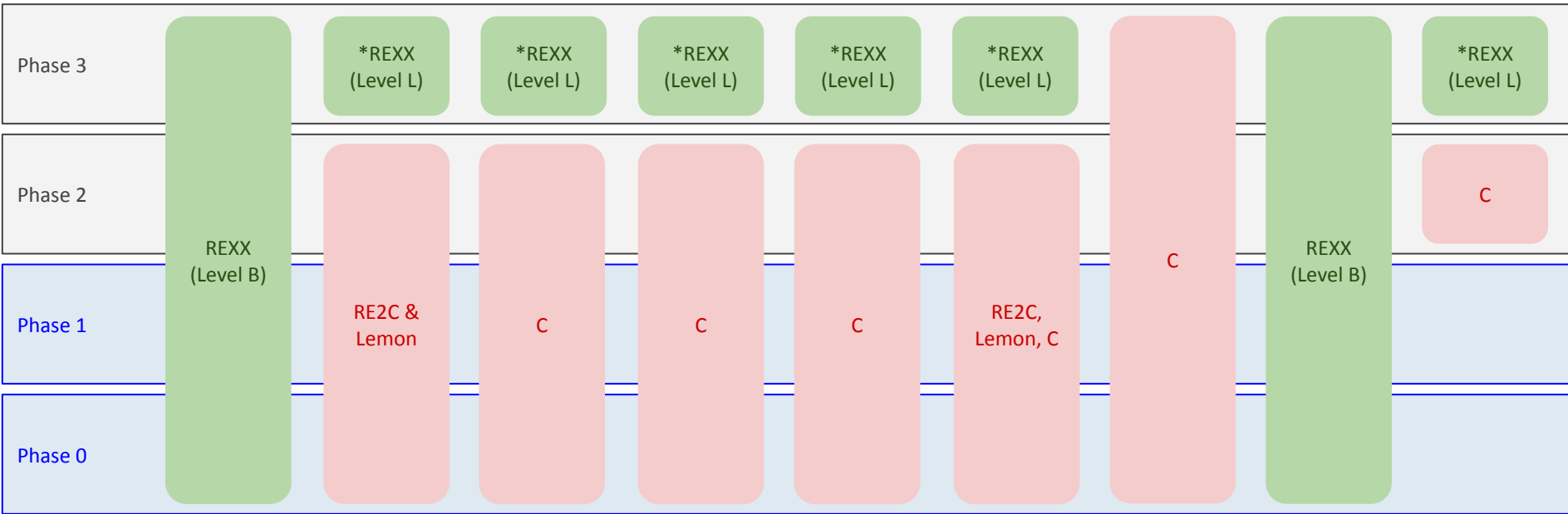
cREXX Architecture

Phases 0 to 2





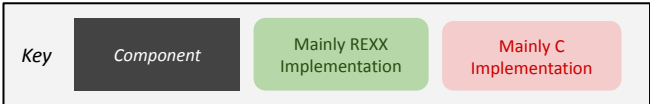


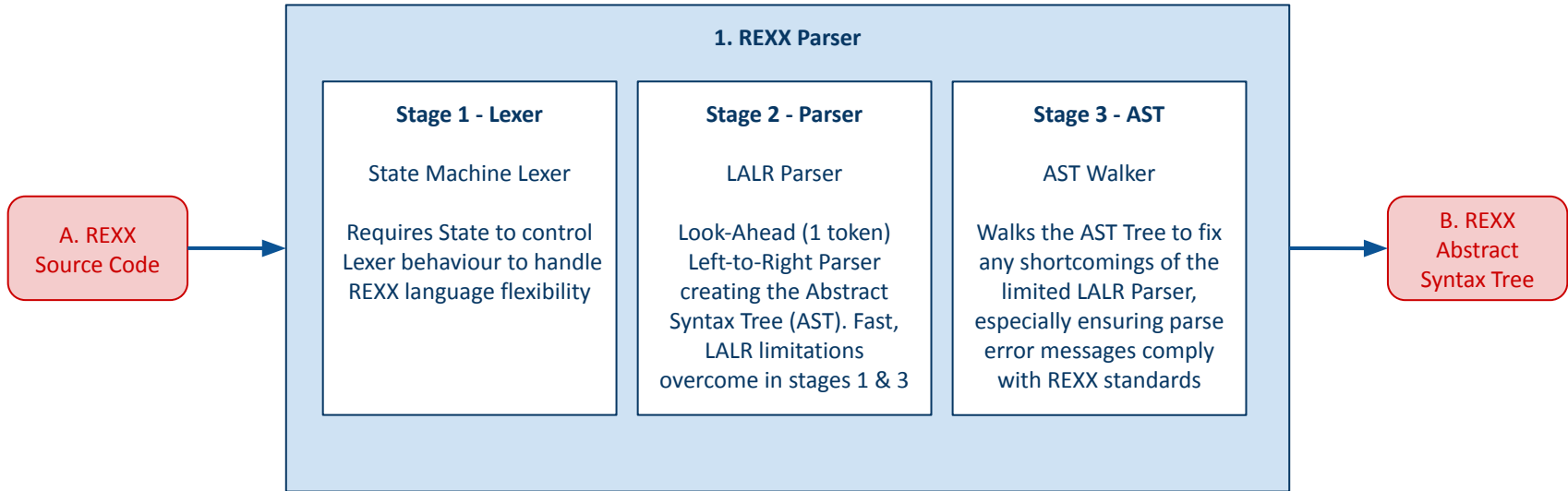


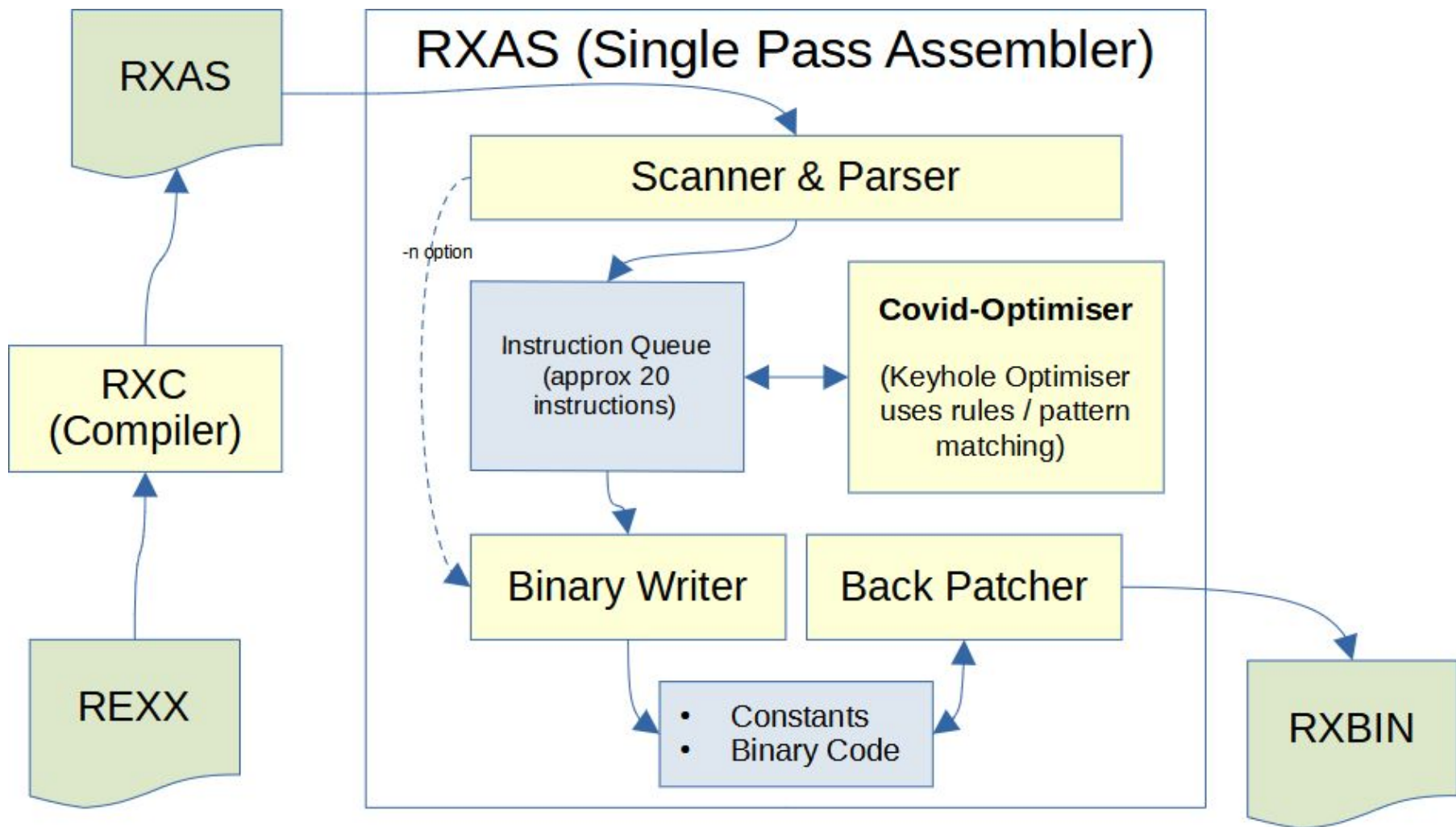
- 0. cREXX End2end Controller
- 1. REXX Parser
- 2. REXX Validator
- 3. REXX Optimiser
- 4. REXX Assembler Generator
- 5. REXX Assembler
- 6. REXX Bytecode Interpreter
- 7. REXX Runtime Library
- 8. REXX LLVM IR Generator

* REXX Level L provides the required:

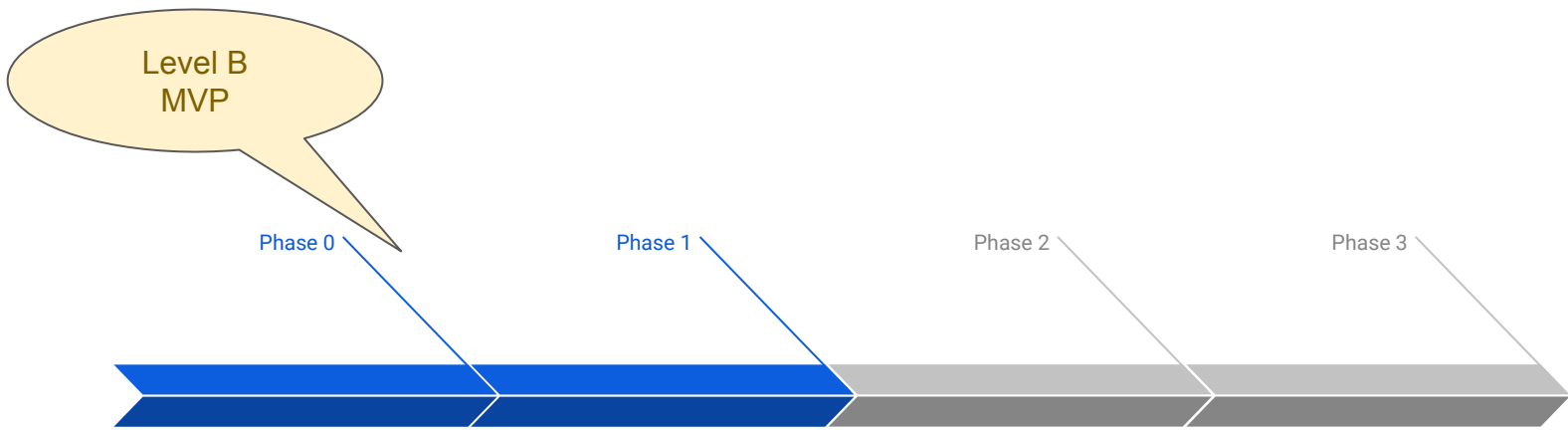
1. Extended PARSE to handle PEG Grammars
2. Native support of Language Engineering data structures (ASTs and Symbol Tables)







cREXX Level B MVP



Level B MVP

Phase 0

Phase 1

Phase 2

Phase 3

Proof of Concept

Goal: Sustainability

Prove architectural concepts and the ability for the project to deliver by creating a modern REXX implementation

Classic REXX

Goal: Standards compliancy

Formalise the implementation by creating a high quality, stable, performant and compliant Classic REXX

Native Performance

Goal: Native Binaries

Integrate to the LLVM backend to allow optimised native binaries for multiple target operating systems

REXX Modernisation

Goal: Contemporary REXX

Re-imagine REXX for new users and workloads, and with contemporary language features

cREXX Level B MVP

Implemented

1. Statically typed language
2. REXX assembler (rxas) based
3. Compiler, Assembler, Interpreter, Debugger (WIP in REXX)
4. Windows, Mac, Linux, VM/370CE + all good C90 targets
5. Metadata (debugging, linking, introspection, interfacing)
6. UTF
7. PROCEDURE, IF, THEN, UNTIL, WHILE, FOREVER, LEAVE, ITERATE, CALL, ARG, SAY, LOOP
8. ASSEMBLER (for low level functionality)
9. Runtime library including runtime “exits” (WIP)
10. Libraries (rxbin)
11. Libraries as “C-Arrays” and linking to standalone native exe’s
12. NAMESPACES and IMPORTing
13. EXPOSE (Static Scoping)
14. EXPOSE across source files
15. Line Comments

To Complete

1. PARSE
2. SELECT
3. Objects
4. Arrays
5. STEM Object (Implemented in REXX)
6. File IO
7. ADDRESS
8. Native Function Calling
9. Inlining
10. Level B “System” Library

Will not include

1. Variable Pool (Level C)
2. SAA Interface (Level B) - To be implemented in REXX
3. LLVM
4. Full Runtime Library (Level C & G)
5. Math[s]

How to Help?



How to Help?

- Github - <https://github.com/adesutherland/CREXX>
- Contact myself or René
- Fortnightly Evening Zoom meetings

- **Code - Test - Use - Feedback - or just lurk!**

Thanks to ...

- **René Jansen** - Our PM; for all his encouragement and work on the built in functions
 - **Peter Jacob** - Our microcode engineer!
 - **Mike Großmann** - For rolling up his sleeves when needed
 - **Michael Beer, Bob Bolch** and everyone else who comes to our project meetings when they should be having a beer!
-

Adrian Sutherland

- Journeyman Architect
- Keeps “hands-on” through numerous projects, from Raspberry PI toys and Domain Specific Languages to open architectural papers and other assets.

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Questions

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Typical Bytecode Optimisation

1. Threaded code
2. Super-instructions / inlining
3. Top-of-stack in a register
4. Scheduling the dispatch of the next VM instruction

In all about 2x faster than classic bytecode

We should be aiming for performance of only 2-5 times slower than native code

NOTE - We could be talking about any language ...

```
char code[] = {
    ICONST_1, ICONST_2,
    IADD, ...
}
char *pc = code;

/* dispatch loop */
while(true) {
    switch(*pc++) {
        case ICONST_1: *++sp = 1; break;
        case ICONST_2: *++sp = 2; break;
        case IADD:
            sp[-1] += *sp; --sp; break;
        ...
    }
}
```

Pure Bytecode

```
void *code[] = {
    &&ICONST_1, &&ICONST_2,
    &&IADD, ...
}
void **pc = code;

/* implementations */
goto **pc);

ICONST_1: pc++; *++sp = 1; goto **pc);
ICONST_2: pc++; *++sp = 2; goto **pc);
IADD:
    pc++; sp[-1] += *sp; --sp; goto **pc);
...
```

Threaded Interpreter

```
/* SIMPLE */  
A = 10  
B = 5  
SAY A + B
```

REXX Assembler

This is where optimisations become REXX specific ...

BREXX

- Stack Based
- Leaves work to the interpreter

CREXX

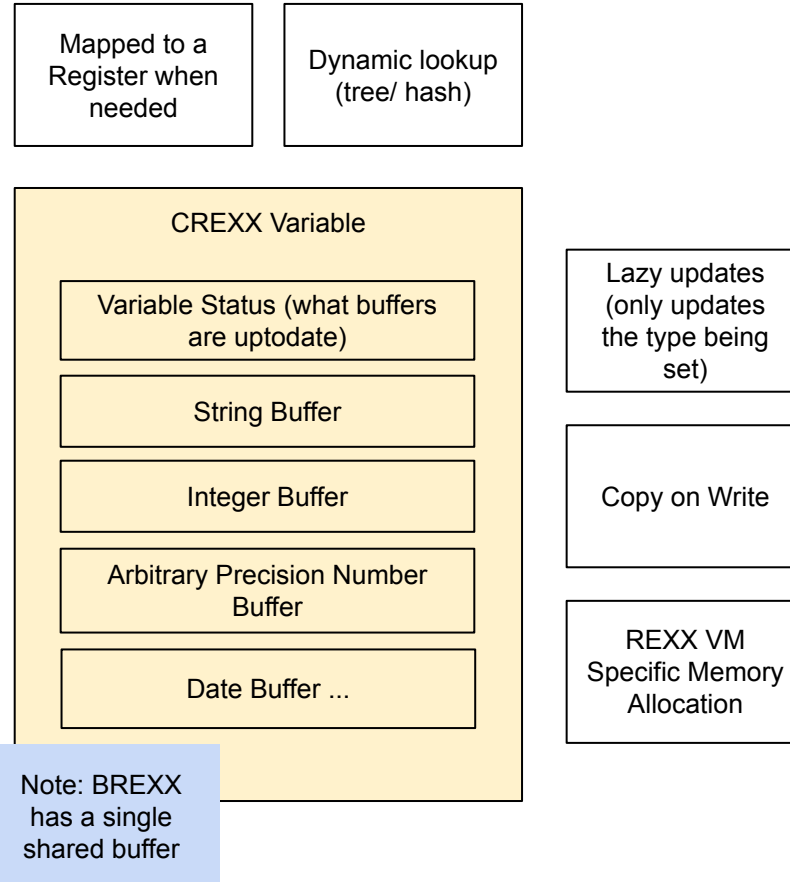
- Register Based
- Trying to handle REXXisms at the low level

<pre>NEWCLAUSE CREATE "A" PUSH 10 COPY NEWCLAUSE CREATE "B" PUSH 5 COPY NEWCLAUSE PUSHTMP LOAD "A" LOAD "B" ADD SAY NEWCLAUSE IEXIT</pre>	<pre>.def main: locals=3 {r1="A", r2="B"} ILOAD r1,10 ILOAD r2,5 IADD r3,r1,r2 ISAY r3 HALT</pre>
BREXX	CREXX

We need to get this right for LLVM ...

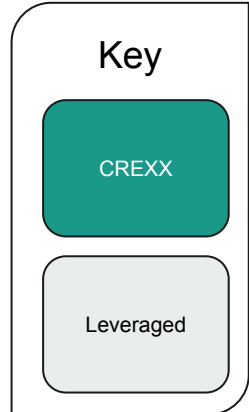
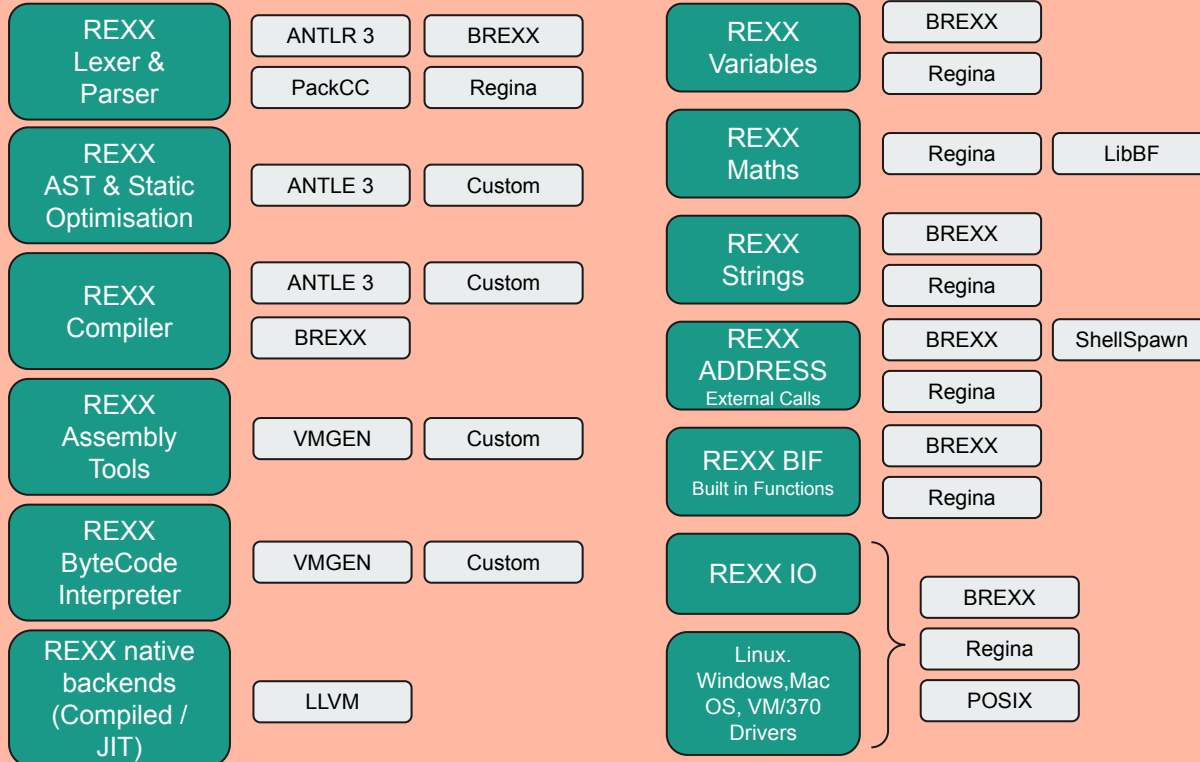
REXX Variable Types

1. REXX is typeless ... and more than that conceptually all variables are strings
2. REXX stems provide a flexible and arbitrary index scheme
3. VALUE(), INTERPRET(), and REXXSAA/EXECOMM all require dynamic variable name resolution
4. Performance requires compile time resolution of variable names and types, wherever possible



CREXX

C Library



High Level Components & Implementation Leverageable Tools